

# A New Coding Scheme for Discrete Memoryless MACs with Common Rate-Limited Feedback

Selma Belhadj Amor

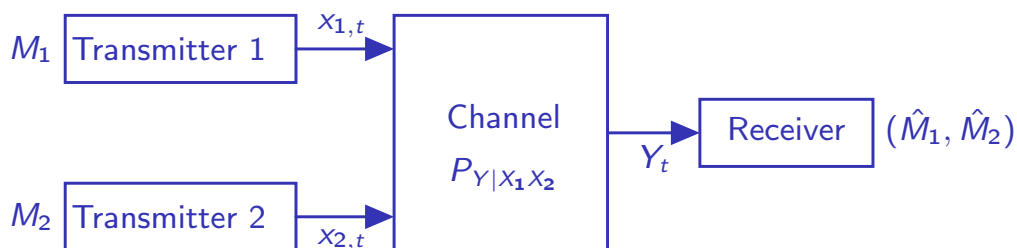
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## Discrete Memoryless Multi-Access Channel

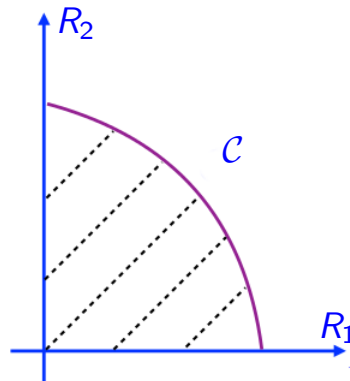


- Finite inputs and output alphabets  $\mathcal{X}_1, \mathcal{X}_2, \mathcal{Y}$
- Memoryless channel  $P_{Y|X_1X_2}$
- Messages  $M_1$  and  $M_2$  independent
- $p(\text{error}) = \Pr((M_1, M_2) \neq (\hat{M}_1, \hat{M}_2))$

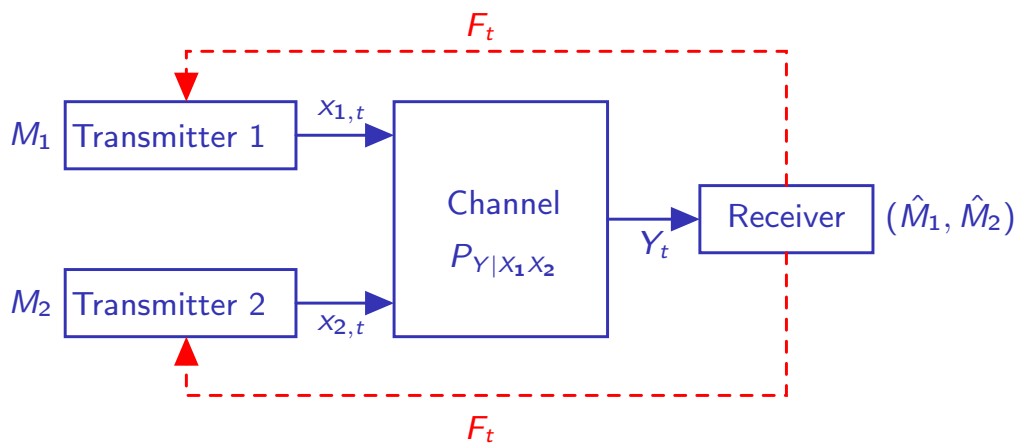
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# Capacity Region

- $M_1 \sim \mathcal{U}\{1, \dots, \lfloor 2^{nR_1} \rfloor\}$  and  $M_2 \sim \mathcal{U}\{1, \dots, \lfloor 2^{nR_2} \rfloor\}$
- $n$ : blocklength
- $R_1, R_2 \geq 0$ : rates of communications
- **Capacity-region  $\mathcal{C}$** : Pairs  $(R_1, R_2)$  s.t.  $p(\text{error})$  can be made arbitrarily small.

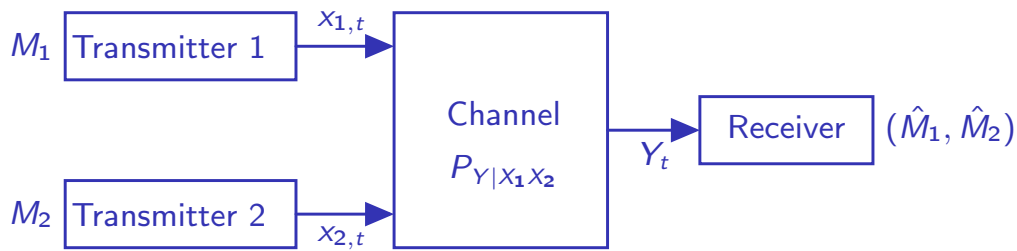


# Rate-Limited Feedback Pipes ( $R_{fb} \geq 0$ )



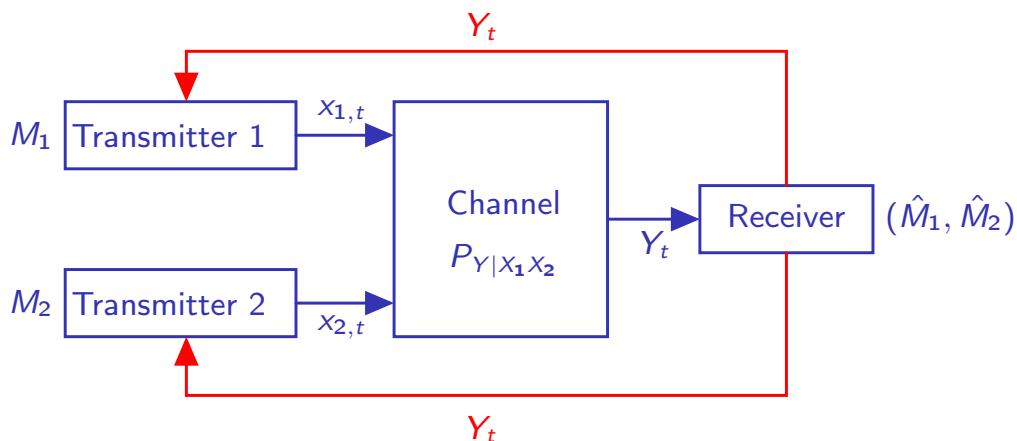
- Feedback signals:  $F_t = \psi_t^{(n)}(Y_1, \dots, Y_t) \in \mathcal{F}_t$ .
- Both transmitters receive feedback signals perfectly whenever
 
$$|\mathcal{F}_1| \times \dots \times |\mathcal{F}_n| \leq 2^{nR_{fb}}.$$
- Inputs:  $X_{i,t} = \varphi_{i,t}^{(n)}(M_i, F_1, \dots, F_{t-1})$

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- Nofeedback:  $R_{\text{fb}} = 0$
- Perfect feedback:  $R_{\text{fb}} \geq \log_2 |\mathcal{Y}|$

# Nofeedback Capacity Region ( $R_{fb} = 0$ )

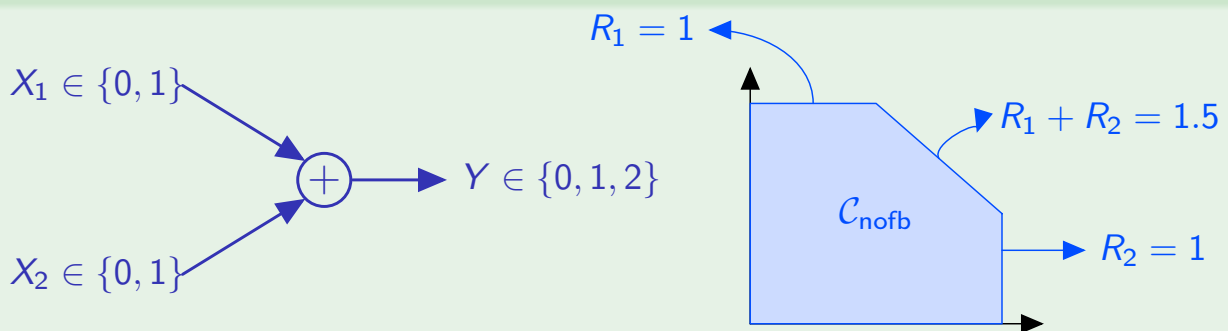
Nofeedback capacity region  $C_{nofb}$  (Ahlsvede'71 and Liao'72)

All nonnegative rate-pairs  $(R_1, R_2)$ :

$$\begin{aligned} R_1 &\leq I(X_1; Y|X_2), \\ R_2 &\leq I(X_2; Y|X_1), \\ R_1 + R_2 &\leq I(X_1X_2; Y), \end{aligned}$$

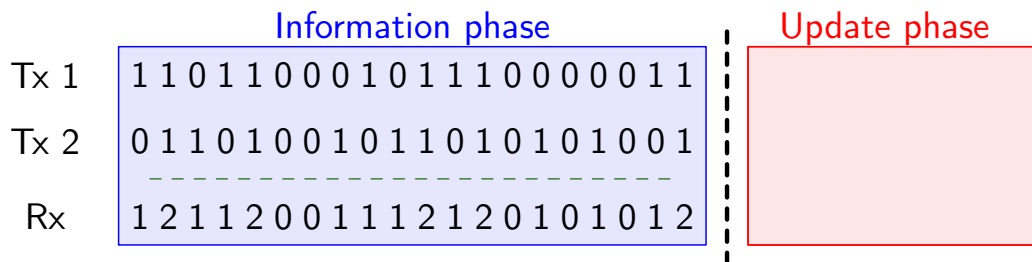
for some  $(X_1, X_2) \sim P_{X_1}P_{X_2}$ .

## Example



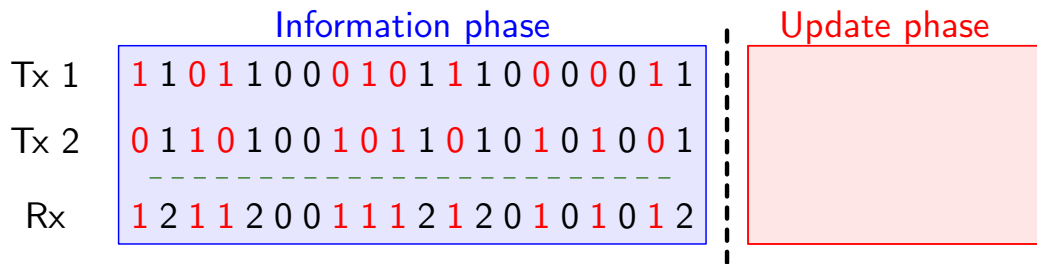
4 / 16

## Feedback Increases Capacity – Gaarder & Wolf 1975



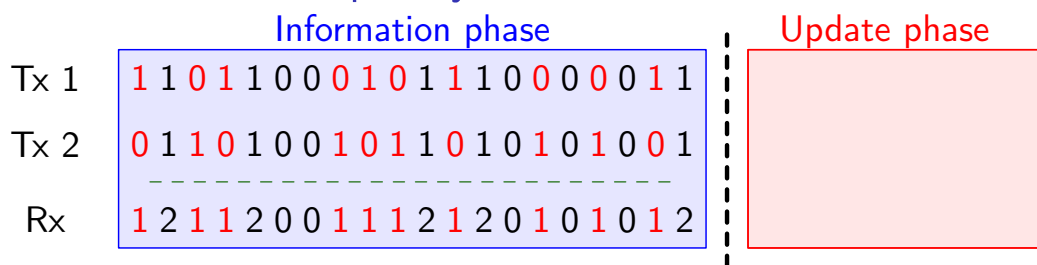
- Each Tx sends  $K$  uncoded bits

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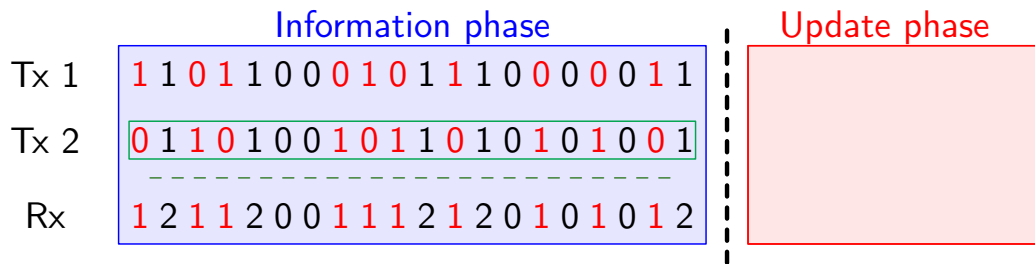
- Each Tx sends  $K$  uncoded bits
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- On average  $\frac{K}{2}$  collisions!

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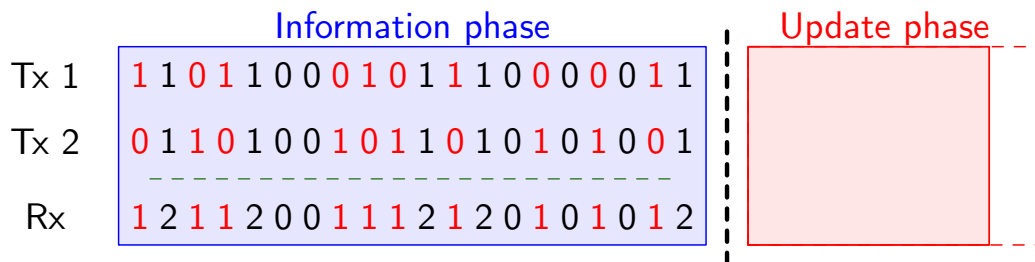
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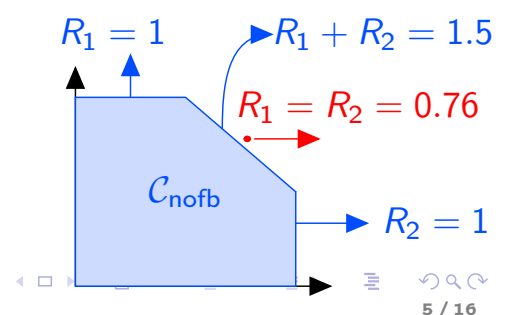
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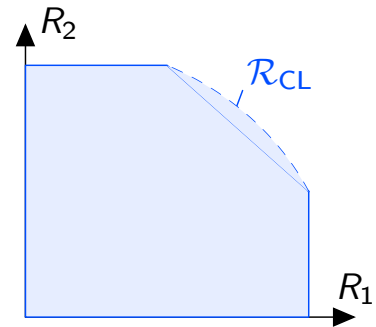


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- Txs cooperate and use 3 input symbols over a noiseless channel at rate  $\log_2 3$



# Cover-Leung Achievable Region $\mathcal{R}_{CL}$



- Cover & Leung 1981 extended G-W scheme
- Interleaving 2 phases of G-W scheme

## Cover-Leung Achievable Region $\mathcal{R}_{CL}$

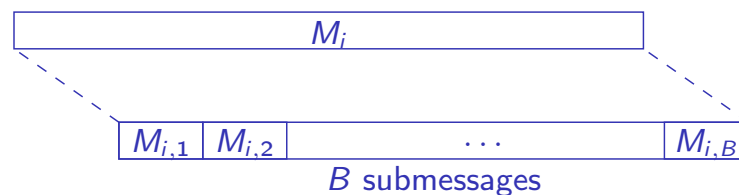
All nonnegative rate-pairs  $(R_1, R_2)$  satisfying

$$\begin{aligned} R_1 &\leq I(X_1; Y|X_2 W), \\ R_2 &\leq I(X_2; Y|X_1 W), \\ R_1 + R_2 &\leq I(X_1 X_2; Y), \end{aligned}$$

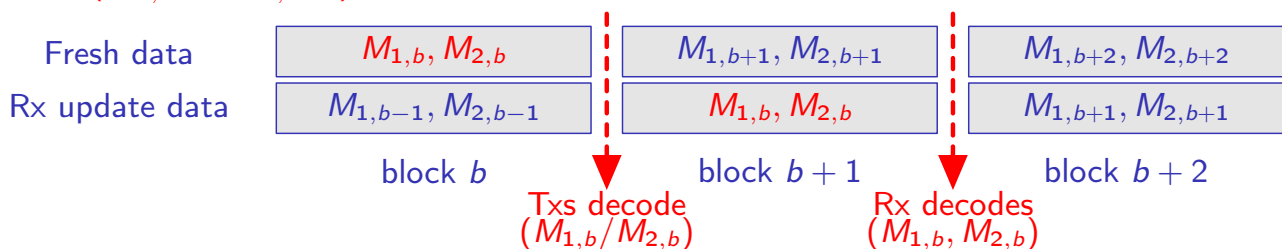
for some  $P_W P_{X_1|W} P_{X_2|W}$  s.t.  $|\mathcal{W}| \leq \min\{|\mathcal{X}_1| |\mathcal{X}_2| + 1, |\mathcal{Y}| + 2\}$ .

## Cover-Leung Coding Scheme

- Each message  $M_i$  is split into  $B$  submessages:



- Block-Markov coding over  $B + 1$  blocks
- In Block  $b$ , send **fresh data**  $(M_{1,b}, M_{2,b})$  superposed on update info about  $(M_{1,b-1}, M_{2,b-1})$ .



- After transmission of Block  $b$ :
  - ▶ Tx 1 (2) decodes  $M_{2,b}$  ( $M_{1,b}$ )
  - ▶ Rx decodes  $(M_{1,b-1}, M_{2,b-1})$





# Previous Results with Common Rate-Limited Feedback

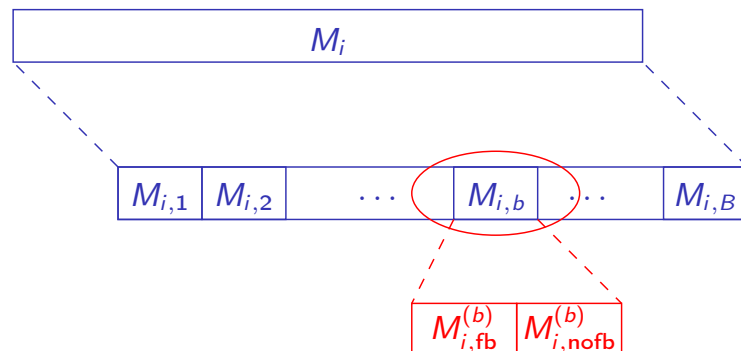
- Shaviv & Steinberg (2008)
  - ▶ Carleial's extension of C-L scheme is used over the "forward channel"
    - ★ If in C-L scheme feedback is 'imperfect' → TxS can still understand useful update info; TxS can decode each other message or some part of it
  - ▶ Heegard-Berger source-coding on the feedback links
- No feedback  $\Rightarrow \mathcal{R}_{SS} = C_{\text{nofb}}$
- Perfect feedback  $\Rightarrow \mathcal{R}_{SS} = \mathcal{R}_{CL}$

Can we do better?

# New Coding Scheme with Common Rate-Limited Feedback

## Main Idea

- In S-S scheme, replace Carleial/C-L scheme by V-P scheme
- Each message  $M_i$  is split into  $B$  submessages
- Each submessage  $M_i^{(b)}$  is split into  $M_{i,\text{fb}}^{(b)}$  and a  $M_{i,\text{nofb}}^{(b)}$

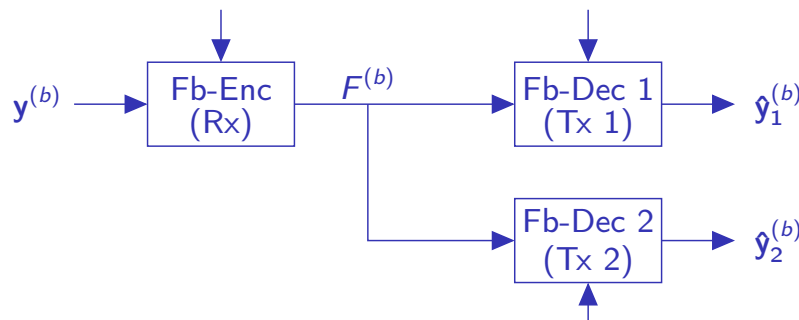


## New Coding Scheme with Common Rate-Limited Feedback

- Superposition & block-Markov coding over  $B + 2$  blocks
- $M_{i,\text{fb}}^{(b)}$  is sent using a **feedback code** (V-P perfect feedback scheme)
- $M_{i,\text{nofb}}^{(b)}$  is sent using a **nofeedback scheme**
  - 1  $(M_{1,\text{fb}}^{(b)}, M_{2,\text{fb}}^{(b)})$  are sent through  $(U_1, U_2)$ -codewords in block  $b$
  - 2 TxS send Tx-side update info through  $(V_1, V_2)$ -codewords in block  $b + 1$ 
    - ★ **correlated info can be sent more efficiently than indep. info**
    - ★ conditions on  $P_{V_1 V_2}$  to ensure stationarity
  - 3 After block  $b + 1$ , TxS decode  $(M_{1,\text{fb}}^{(b)}, M_{2,\text{fb}}^{(b)})$  using feedback
  - 4 TxS jointly send common Rx-side update info about  $(M_{1,\text{fb}}^{(b)}, M_{2,\text{fb}}^{(b)})$  through  $W$ -codewords in block  $b + 2$
  - 5 After block  $b + 2$ , Rx decodes  $(M_{1,\text{fb}}^{(b)}, M_{2,\text{fb}}^{(b)})$ , then decodes  $(M_{1,\text{nofb}}^{(b)}, M_{2,\text{nofb}}^{(b)})$ .

## New Coding Scheme with Common Rate-Limited Feedback

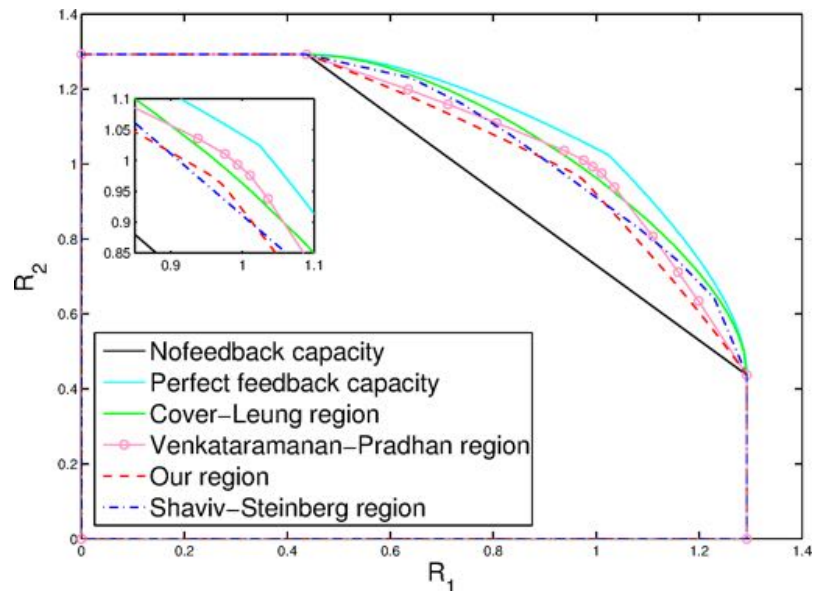
- Feedback links are used at the end of each block to send **compression information** about channel outputs
- Feedback encoding = **Heegard-Berger'85 source-coding**:





# Achievable Region

- No feedback:  $\mathcal{R}_{NEW} = \mathcal{C}_{nofb}$
- $\mathcal{R}_{NEW} \supseteq \mathcal{R}_{SS}$  and inclusion can be strict:
  - ▶ Perfect feedback:  $\mathcal{R}_{NEW} = \mathcal{R}_{VP} \supset \mathcal{R}_{CL} = \mathcal{R}_{SS}$
  - ▶ Specific (Gaussian) choice of RVs for 2-user Gaussian MAC  $\rightarrow$  New scheme achieves a **strictly larger sumrate** than S-S scheme



Achievable regions for the Gaussian MAC for  $P/\sigma^2 = 5$  and  $R_{fb} = 2$